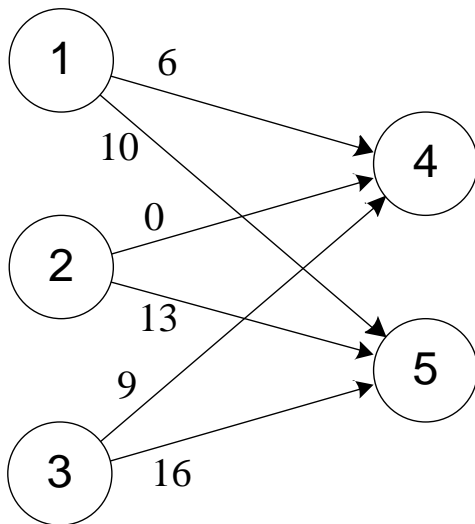


Topics

1. Introduction
2. Facility location
3. Freight transport
 - Exam 1 (take home)
- 4. Network models**
5. Routing
 - Exam 2 (take home)
6. Warehousing
 - Final exam (in class)

Graph Representations

- Complete bipartite directed (or digraph):
 - Suppliers to multiple DCs, single mode of transport



C: 1 2
 ---:-----
 1: 6 10
 2: 0 13
 3: 9 16

Interlevel matrix

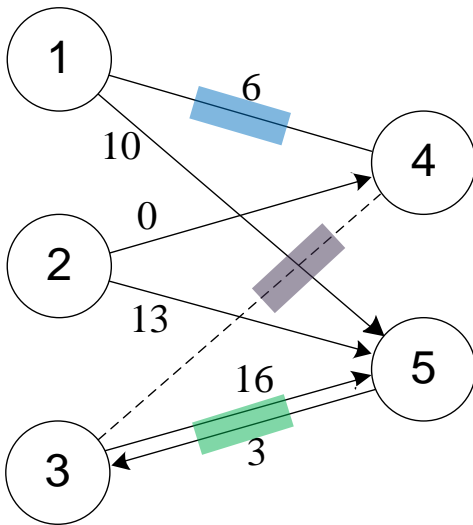
W =

0	0	0	6	10
0	0	0	NaN	13
0	0	0	9	16
0	0	0	0	0
0	0	0	0	0

Weighted adjacency matrix

Graph Representations

- Bipartite:
 - One- or two-way connections between nodes in two groups



$$W = \begin{matrix} & & & \begin{matrix} \textcircled{6} & 10 \\ \text{NaN} & 13 \\ \textcircled{0} & \textcircled{16} \\ 0 & 0 \end{matrix} \\ \begin{matrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \\ \textcircled{6} & 0 & 0 \\ 0 & 0 & \textcircled{3} \end{matrix} & & & \begin{matrix} 0 & 0 \\ 0 & 0 \\ 0 & 0 \end{matrix} \end{matrix}$$

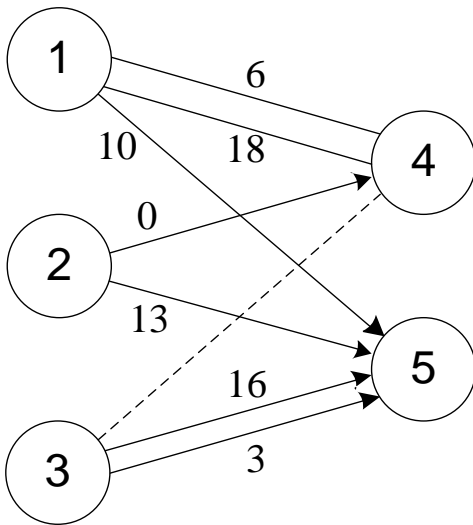
$$IJC = \begin{matrix} \begin{matrix} 4 & 1 & 6 \\ 5 & 3 & 3 \\ 1 & 4 & 6 \\ 2 & 4 & 0 \\ 1 & 5 & 10 \\ 2 & 5 & 13 \\ 3 & 5 & 16 \end{matrix} & \left. \vphantom{\begin{matrix} 4 \\ 5 \\ 1 \\ 2 \\ 1 \\ 2 \\ 3 \end{matrix}} \right\} \end{matrix}$$

Arc list matrix



Graph Representations

- Multigraph:
 - Multiple connections, multiple modes of transport



```

IJC = 1    -4    6
      1    -4   18
      1     5   10
      2     4    0
      2     5   13
      3     5   16
      3     5    3
    
```

```

no_W =

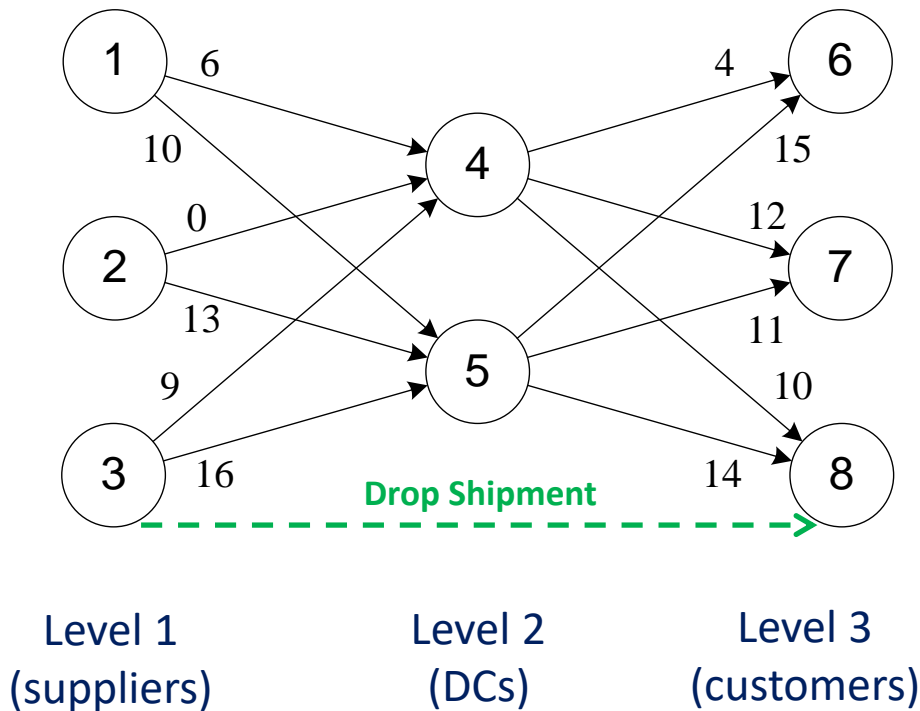
0    0    0    24    10
0    0    0    NaN   13
0    0    0     0   19
24   0    0     0    0
0    0    0     0    0


```

Can't represent using adjacency matrix

Graph Representations

- Complete multipartite directed:
 - Typical supply chain (no drop shipments)

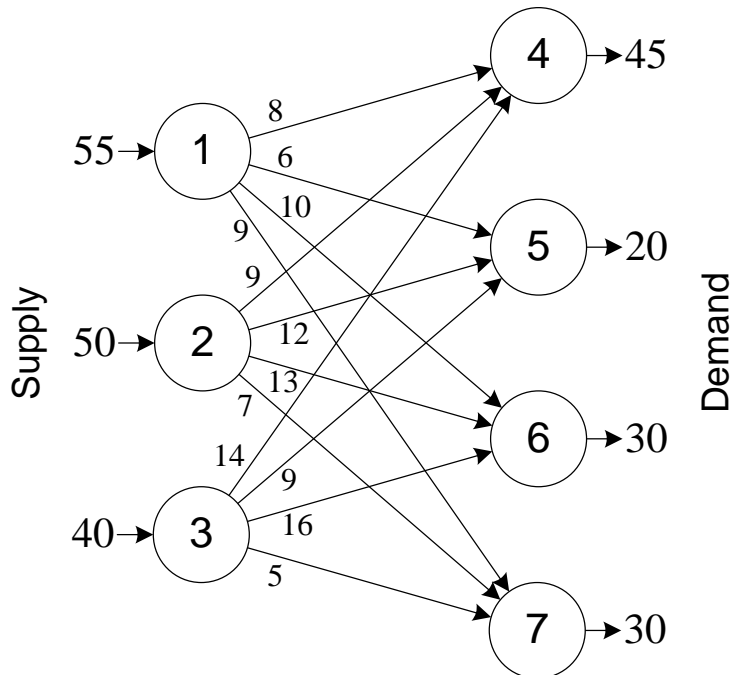


C12:	1	2	C23:	1	2	3
1:	6	10	1:	4	12	10
2:	0	13	2:	15	11	14
3:	9	16				

W:	1	2	3	4	5	6	7	8
1:	0	0	0	6	10	0	0	0
2:	0	0	0		13	0	0	0
3:	0	0	0	9	16	0	0	0
4:	0	0	0	0	0	4	12	10
5:	0	0	0	0	0	15	11	14
6:	0	0	0	0	0	0	0	0
7:	0	0	0	0	0	0	0	0
8:	0	0	0	0	0	0	0	0

Transportation Problem

- Satisfy node demand from supply nodes
 - Can be used for allocation in ALA when NFs have capacity constraints
 - Min cost/distance allocation = infinite supply at each node



Trans	4	5	6	7	Supply
1	8	6	10	9	55
2	9	12	13	7	50
3	14	9	16	5	40
Demand	45	20	30	30	

Greedy Solution Procedure

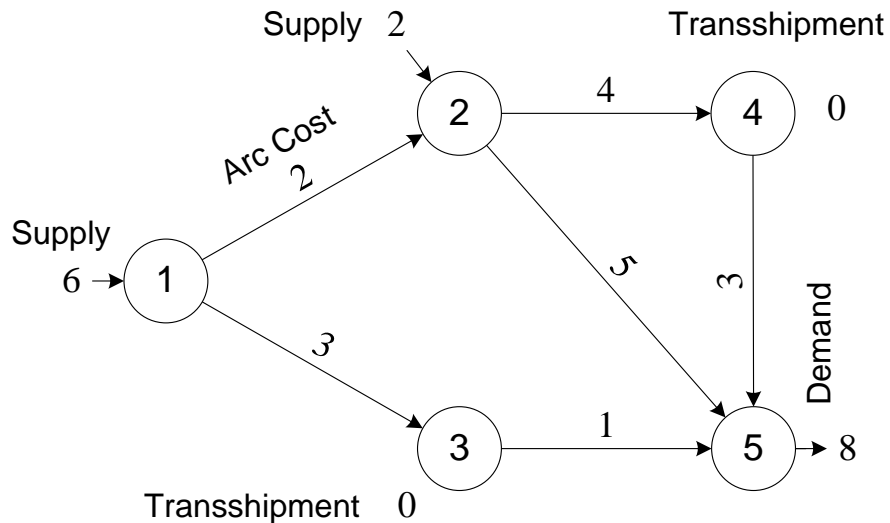
- Procedure for transportation problem: *Continue to select lowest cost supply until all demand is satisfied*
 - Fast, but not always optimal for transportation problem
 - Dijkstra's shortest path and simplex method for LP are optimal greedy procedures

Trans	4	5	6	7	Supply
1	8	6	10	9	55 -20 = 35 -35 = 0
2	9	12	13	7	50 -10 = 40 -30 = 10
3	14	9	16	5	40 -30 = 10
Demand	45	20	30	30	
	10 0	0	0	0	

$$TC = 5(30) + 6(20) + 8(35) + 9(10) + 13(30) = 1,030 \quad (\text{vs } 970 \text{ optimal})$$

Min Cost Network Flow (MCNF) Problem

- Most general network problem, can solve using any type of graph representation



$$s_i = \text{net supply of node } i$$

$$= \begin{cases} > 0, & \text{supply node} \\ < 0, & \text{demand node} \\ = 0, & \text{transshipment node} \end{cases}$$

$$\text{Arc cost: } \mathbf{c} = [2 \ 3 \ 4 \ 5 \ 1 \ 3]'$$

$$\text{Net node supply: } \mathbf{s} = [6 \ 2 \ 0 \ 0 \ -8]'$$

$$\text{Incidence Matrix: } \mathbf{A} = \begin{bmatrix} 1 & 1 & 0 & 0 & 0 & 0 \\ -1 & 0 & 1 & 1 & 0 & 0 \\ 0 & -1 & 0 & 0 & 1 & 0 \\ 0 & 0 & -1 & 0 & 0 & 1 \\ 0 & 0 & 0 & -1 & -1 & -1 \end{bmatrix}$$

MCNF:	lhs	C	C	C	C	C	C	rhs
Min:		2	3	4	5	1	3	
1:	6	1	1	0	0	0	0	6
2:	2	-1	0	1	1	0	0	2
3:	0	0	-1	0	0	1	0	0
4:	0	0	0	-1	0	0	1	0
lb:		0	0	0	0	0	0	
ub:		Inf	Inf	Inf	Inf	Inf	Inf	

Row for node 5 is redundant



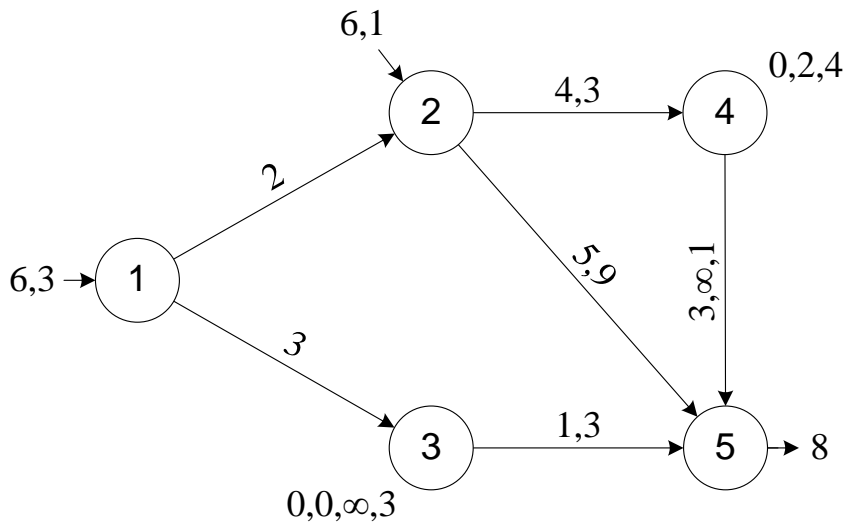
$$\text{MCNF: } \max \mathbf{c}'\mathbf{x}$$

$$\text{s.t. } \mathbf{Ax} = \mathbf{s}$$

$$\mathbf{x} \geq 0$$

MCNF with Arc/Node Bounds and Node Costs

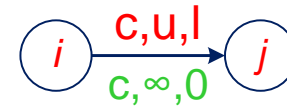
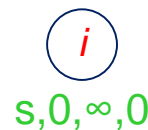
- Bounds on arcs/nodes can represent capacity constraints in a logistic network
- Node cost can represent production cost or intersection delay



s_i = net supply of node i

$$= \begin{cases} > 0, & \text{supply node} \\ < 0, & \text{demand node} \\ = 0, & \text{transshipment node} \end{cases}$$

s, nc, nu, nl



IJCUL: 1 2 3 4 5

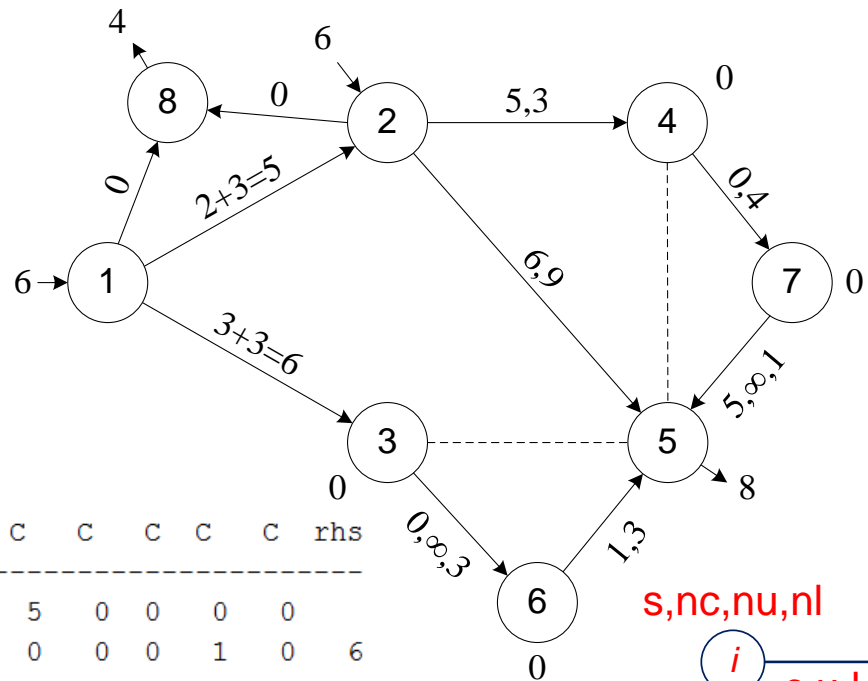
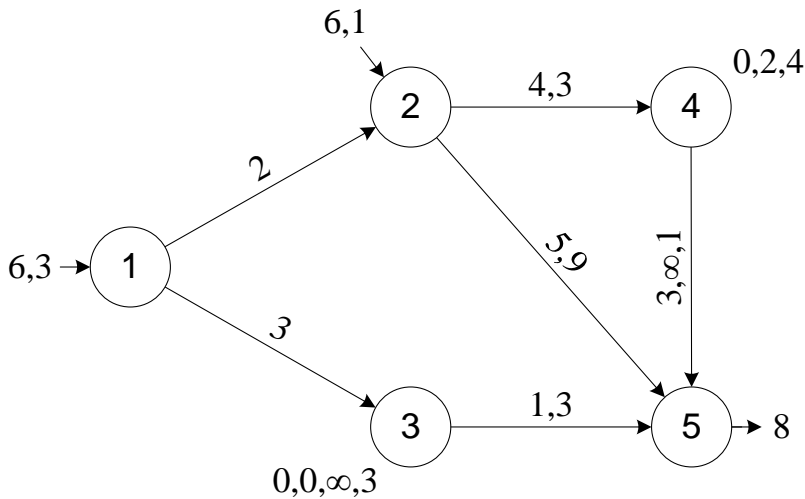
1:	1	2	2	Inf	0
2:	1	3	3	Inf	0
3:	2	4	4	3	0
4:	2	5	5	9	0
5:	3	5	1	3	0
6:	4	5	3	Inf	1

SCUL: 1 2 3 4

1:	6	3	Inf	0
2:	6	1	Inf	0
3:	0	0	Inf	3
4:	0	2	4	0
5:	-8	0	Inf	0

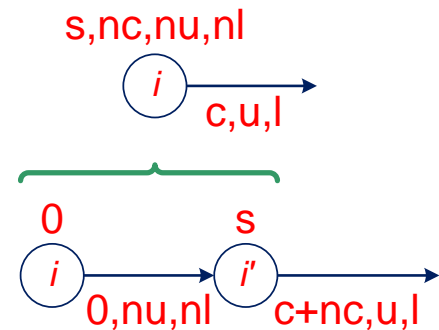
Expanded-Node Formulation of MCNF

- Node cost/constraints converted to arc cost/constraints
 - Dummy node (8) added so that supply = demand



MCNF: lhs C C C C C C C C C C C rhs

Min:		5	6	5	6	1	5	0	0	0	0	
1:	6	1	1	0	0	0	0	0	0	1	0	6
2:	6	-1	0	1	1	0	0	0	0	0	1	6
3:	0	0	-1	0	0	0	0	1	0	0	0	0
4:	0	0	0	-1	0	0	0	0	1	0	0	0
5:	-8	0	0	0	-1	-1	-1	0	0	0	0	-8
6:	0	0	0	0	0	1	0	-1	0	0	0	0
7:	0	0	0	0	0	0	1	0	-1	0	0	0
lb:		0	0	0	0	0	1	3	0	0	0	
ub:		Inf	Inf	3	9	3	Inf	Inf	4	Inf	Inf	



Solving an MCNF as an LP

- Special procedures more efficient than LP were developed to solve MCNF and Transportation problems
 - e.g., Network simplex algorithm (MCNF)
 - e.g., Hungarian method (Transportation and *Transshipment*)
- Now usually easier to transform into LP since solvers are so good, with MCNF just aiding in formulation of problem:
 - Trans \Rightarrow MCNF \Rightarrow LP
 - Special, very efficient procedures only used for shortest path problem (Dijkstra)

